# CONTRACTING



## $\mathbf{T} = f(\$, \$, \mathbf{I}, \mathbf{mp}, ...)$

- § : law & regulation
- *\$* : financing
- **I**: location
- **m** : technology
- **p**: time period

## **GRAPH** (Basic Terms of Graph Techniques )

#### As a "model" :

Strictly identified elements and well defined ,,bilateral" (pairly) relations among them ...

#### **Elements :**

- particles
- phases / states
- processes

#### **Relations :**

- links, joints
- cause-result interactions
- precedences, sequences

#### •

#### Mathematically :

Structured set of nodes and edges. Edge : related **pair of nodes** ...



**Set of nodes** (N = "node")N = { a, b, c, d, e, f }

**Set of edges** (E = "edge")

 $E = [ \{a,c\}, \{a,e\}, \{b,c\}, \{b,d\}, \\ \{b,e\}, \{c,e\}, \{c,f\}, \{d,f\} ]$ 

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Graph (G = "graph" @ graphics)
G = [N, E]
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#### **Directed Edge** ( *A* = "arc", "arrow" )

Relation between paired nodes  $\{i, j\}$  interpreted in one direction (e.g. from "i" to "j" ) only. ( Order of nodes also indicates direction. E.g.: (i, j), ... (a, e), ...)



 $N = \{ a, b, c, d, e, f \}$ 

 $A = \{ (a,c), (a,e), (b,c), (b,d), \\ (c,b), (c,f), (e,b), (e,c), (f,d) \}$ 

$$G = [N, A]$$

### **Directed Graph**

(frequently referred as "DiGraph")

"A graph with all the edges directed "

(Implicitly: Between two nodes at most one directed edge should be allowed ? ... )

#### **Remark** :

Any "non-directed" edge can be handled as "directed", since any non-directed edge can be substituted with an opposed pair of directed edges between the same two related nodes

 $\{i, j\} = \{(i, j), (j, i)\}$ 

( ... Anyway, why not to allow existence of more directed edges between any related pair of nodes ? ... )



# Weighted Graph

Quantitative characteristics so called ,,weights" are interpreted/assigned by nodes and edges.



# **Basic Terms of Directed Graphs**

#### Source :

A nod being origin of at least one directed edge, but not terminal point of any directed edges.



### Sink :

A nod being terminal point of at least one directed edge, but not origin of any directed edges.



## Path : (P)

Continuous repeatless directed chain (string) of directed edges.

*Identifying them by the sequence of linked nodes. e.g.*: *P*[*i*,*l*] = { *i*, *j*, *k*, *l* }



Loop :

A path with origin and terminal point the same. ,,Self-closing path". e.g.:  $P[i,i] = \{i, j, k, (i)\}$ 



# **Graph - Topologies**

(In relation of nodes and edges, etc. ...)

"comlete"

"bipartite"





"tree"



"connected" "non-connected"



**Structure (''adjacency'') matrix** 





## The "Network"



#### **Network** (*as terminus technicus* ) :

Connected weighted directed graph with a single source and a single sink but with no loops and no negative weights.

#### **Network** ( *as a popular reference* ) :

A graph ... with no any specification or generalization.

## **Network "Problems"** ( the most popular questions )

- Path finding \*

- Integrity (connectivity) analysis
- Loop discovery
- Dominance analysis
- Path-variants
- Longest path / Shortest path \*
- Gravity-point / Center-point
- Maximal flow / Minimal cut \*
- Potentials' problems

#### \* so called "directed problems"

# **Scheduling by Networks**

#### **Network Problem Analogies:**

- The longest path
- The minimum potentials'

(All the elements are relevant, but we are looking for de dominant ones and are trying to predict generated effects of any changes.)

# **Scheduling Techniques using Networks** ( *complementary algorithms, interpretations* )

- PERT<sup>time</sup>
- CPM<sup>time</sup>
- CPM<sup>cost</sup>
- CPM<sup>time+</sup>
- MPM<sup>time</sup>/PDM<sup>time</sup>
- MPM<sup>cost</sup>
- GTM ( General Time Model )